CSI 409 — Fall 2017: Homework #6 Some answers and hints

1. Prove that the language

$$\left\{ a^m b^m c^n \mid n > 2m \ge 0 \right\}$$

is not context-free.

Take $z = a^p b^p c^{2p+1}$ where *p* is the pumping length. By the pumping lemma, there are u, v, w, x, y such that $z = a^p b^p c^{2p+1} = uvwxy$, where *v* and *x* are respectively the left and right pumps. Additionally, $1 \le |vwx| \le p$.

Now if *vwx* is a substring of $a^p b^p$ then pump up. If it's a substring of $b^p c^{2p+1}$ then remove the pumps.

2. Prove that the language

$$\Big\{\,a^ib^j\ \Big|\ (i^2>j\ge 0)\Big\}$$

is not context-free.

Take $z = a^{p+1}b^{p^2+2p}$ where p is the pumping length. By the pumping lemma, there are u, v, w, x, y such that $a^{p+1}b^{p^2+2p} = uvwxy$, where v and x are respectively the left and right pumps. Additionally, $1 \le |vx| \le p$.

Now if *vwx* is a substring of b^{p^2+2p} , then pump up. This increases the number of *b*'s, but the number of *a*'s stays the same.

If *vwx* contains an *a*, then remove the pumps. Let z' = uwy. Now $\#_a(z') \leq p$ and $\#_b(z') = p^2 + 2p - k$ for some $0 \leq k \leq p$. The value of p^2 will always be less than $p^2 + 2p - k$ since $k \leq p$.