

# **Chapter 3: Introduction to SQL**

Database System Concepts, 6<sup>th</sup> Ed.

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# **Chapter 3: Introduction to SQL**

- Overview of the SQL Query Language
- Data Definition
- Basic Query Structure
- Additional Basic Operations
- Set Operations
- Null Values
- Aggregate Functions
- Modification of the Database
- Nested Subqueries



### **History**

- IBM Sequel language developed as part of System R project at the IBM San Jose Research Laboratory
- Renamed Structured Query Language (SQL)
- ANSI and ISO standard SQL:
  - SQL-86
  - SQL-89
  - SQL-92
  - SQL:1999 (language name became Y2K compliant!)
  - SQL:2003
- Commercial systems offer most, if not all, SQL-92 features, plus varying feature sets from later standards and special proprietary features.
  - Not all examples here may work on your particular system.



# **Chapter 3: Introduction to SQL**

Overview of the SQL Query Language

### Data Definition

- domain types
- create table with integrity constraints
- drop table / alter table
- Basic Query Structure
- Additional Basic Operations
- Set Operations
- Null Values
- Aggregate Functions
- Modification of the Database
- Nested Subqueries



# **Domain Types in SQL**

- **char(n).** *Fixed length* character string, with user-specified length *n*.
- **varchar(n).** Variable length character strings, with user-specified maximum length n.
- **int.** Integer (a finite subset of the integers that is machine-dependent).
- **smallint.** Small integer (a machine-dependent subset of the integer domain type).
- numeric(p,d). Fixed point number, with user-specified precision of p digits, with d digits to the right of decimal point.
- real, double precision. Floating point and double-precision floating point numbers, with machine-dependent precision.
- float(n). Floating point number, with user-specified precision of at least n digits.
- More are covered in Chapter 4.



# **Create Table Construct**

An SQL relation is defined using the **create table** command:

create table  $r (A_1 D_1, A_2 D_2, ..., A_n D_n, (integrity-constraint_1),$ 

(integrity-constraint<sub>k</sub>))

- *r* is the name of the relation
- each  $A_i$  is an attribute name in the schema of relation r
- $D_i$  is the data type of values in the domain of attribute  $A_i$

Example:

create table instructor (<br/>IDIDchar(5),namevarchar(20) not null,dept\_namevarchar(20),salarynumeric(8,2))

- insert into instructor values ('10211', 'Smith', 'Biology', 66000);
- **insert into** *instructor* **values** ('10211', null, 'Biology', 66000);



# **Integrity Constraints in Create Table**

- not null
- **primary key**  $(A_1, ..., A_n)$
- **foreign key**  $(A_m, ..., A_n)$  references *r*

Example: Declare *ID* as the primary key for *instructor* 

create table instructor ( ID char(5), name varchar(20) not null, dept\_name varchar(20), salary numeric(8,2), primary key (ID), foreign key (dept\_name) references department)

primary key declaration on an attribute automatically ensures not null



### **Drop and Alter Table Constructs**

drop table

#### alter table

- alter table *r* add *A D* 
  - where A is the name of the attribute to be added to relation
     r and D is the domain of A.
  - All tuples in the relation are assigned *null* as the value for the new attribute.
- alter table *r* drop *A* 
  - where A is the name of an attribute of relation r
  - Dropping of attributes not supported by many databases.



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  - select / from / where / natural join
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# **Basic Query Structure**

A typical SQL query has the form:

**select**  $A_1, A_2, ..., A_n$ from  $r_1, r_2, ..., r_m$ where *P* 

- $A_i$  represents an attribute
- $R_i$  represents a relation
- *P* is a predicate.
- The result of an SQL query is a relation.
- NOTE: SQL names are case insensitive



# **The select Clause**

The **select** clause list the attributes desired in the result of a query

- corresponds to the projection operation of the relational algebra
  - select dept\_name
    from instructor



# **The select Clause**

The **select** clause list the attributes desired in the result of a query

corresponds to the projection operation of the relational algebra

select distinct dept\_name

from instructor



# **The select Clause**

- The **select** clause list the attributes desired in the result of a query
  - corresponds to the projection operation of the relational algebra
    - select distinct dept\_name
      - from instructor
- SQL allows duplicates in relations as well as in query results.
- To force the elimination of duplicates, insert the keyword **distinct** after select.
- Find the names of all departments with instructor, and remove duplicates The keyword **all** specifies that duplicates not be removed.

select all dept\_name from instructor



# The select Clause (Cont.)

An asterisk in the select clause denotes "all attributes"

select \* from instructor

- The select clause can contain arithmetic expressions involving the operation, +, -, \*, and /, and operating on constants or attributes of tuples.
  - select ID, name, salary/12
     from instructor



# The where Clause

The **where** clause specifies conditions that the result must satisfy

- Corresponds to the *selection predicate* of the relational algebra.
- To find all instructors in Comp. Sci. dept with salary > 80000

select name
from instructor
where dept\_name = 'Comp. Sci.' and salary > 80000

- Comparison results can be combined using the logical connectives and, or, and not.
  - Comparisons can be applied to results of arithmetic expressions.



# **The from Clause**

The **from** clause lists the relations involved in the query

- Corresponds to the *Cartesian product* operation of the relational algebra.
- Find the Cartesian product *instructor X teaches*

select \*
from instructor, teaches

- generates every possible instructor teaches pair, with all attributes from both relations.
- Cartesian product not very useful directly, but useful combined with where-clause condition (selection operation in relational algebra).



### **Joins**

For all instructors who have taught courses, find their names and the course ID of the courses they taught.

select name, course\_id
from instructor, teaches
where instructor.ID = teaches.ID

#### instructor

teaches

10101         Srinivasan         Comp. Sci.         65000         10101         CS-101         1         Fall           12121         Wu         Finance         90000         10101         CS-315         1         Spring           15151         Mozart         Music         40000         10101         CS-347         1         Fall	year	semester	sec_id	course_id	ID	salary	dept_name	name	ID
22222         Einstein         Physics         95000         12121         FIN-201         1         Spring           32343         El Said         History         60000         15151         MU-199         1         Spring           22222         PHY-101         1         Fall	2009 2010 2009 2010 2010 2010 2009	Spring Fall Spring	1 1 1 1	CS-315 CS-347 FIN-201 MU-199	10101 10101 12121	65000 90000 40000 95000	Comp. Sci. Finance Music Physics	Srinivasan Wu Mozart Einstein	10101 12121 15151 22222



# **Try Writing Some Queries in SQL**

Can you write the following queries in SQL?

• Find the IDs of the courses taught by Srinivasan.

• Find all the departments that offered courses in the Fall of 2009.

#### instructor

ID	name	dept_name	salary
10101	Srinivasan	Comp. Sci.	65000
12121	Wu	Finance	90000
15151	Mozart	Music	40000
22222	Einstein	Physics	95000
32343	El Said	History	60000
00151		DI	07000

#### teaches

ID	course_id	sec_id	semester	year
10101	CS-101	1	Fall	2009
10101	CS-315	1	Spring	2010
10101	CS-347	1	Fall	2009
12121	FIN-201	1	Spring	2010
15151	MU-199	1	Spring	2010
22222	PHY-101	1	Fall	2009

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# **Try Writing Some Queries in SQL**

Can you write the following queries in SQL?

Find the IDs of the courses taught by Srinivasan. **select distinct** course\_id **from** teachers, instructor **where** instructor.name = `Srinivasan' **and** instructor.ID = teachers.ID

• Find all the departments that offered courses in the Fall of 2009.

#### instructor

ID	name	dept_name	salary	
10101	Srinivasan	Comp. Sci.	65000	
12121	Wu	Finance	90000	
15151	Mozart	Music	40000	
22222	Einstein	Physics	95000	
32343	El Said	History	60000	
00154		DI	07000	

#### teaches

ID	course_id	sec_id	semester	year
10101	CS-101	1	Fall	2009
10101	CS-315	1	Spring	2010
10101	CS-347	1	Fall	2009
12121	FIN-201	1	Spring	2010
15151	MU-199	1	Spring	2010
22222	PHY-101	1	Fall	2009

16



# **Try Writing Some Queries in SQL**

Can you write the following queries in SQL?

Find the IDs of the courses taught by Srinivasan. **select distinct** course\_id **from** teachers, instructor **where** instructor.name = `Srinivasan' **and** instructor.ID = teachers.ID

Find all the departments that offered courses in the Fall of 2009.
 select distinct dept\_name

from instructor, teachers

where semester = `Fall' and year = 2009 and teaches.ID = instructor.ID

ID	name	dept_name	salary	
10101	Srinivasan	Comp. Sci.	65000	
12121	Wu	Finance	90000	
15151	Mozart	Music	40000	
22222	Einstein	Physics	95000	
32343	El Said	History	60000	
00154		DI	07000	

#### instructor

#### teaches

ID	course_id	sec_id	semester	year
10101	CS-101	1	Fall	2009
10101	CS-315	1	Spring	2010
10101	CS-347	1	Fall	2009
12121	FIN-201	1	Spring	2010
15151	MU-199	1	Spring	2010
22222	PHY-101	1	Fall	2009



# **Natural Join**

Natural join matches tuples with the same values for all common attributes, and retains only one copy of each common column

#### select \*

from instructor natural join teaches;

ID	name	dept_name	salary	course_id	sec_id	semester	year
10101	Srinivasan	Comp. Sci.	65000	CS-101	1	Fall	2009
10101	Srinivasan	Comp. Sci.	65000	CS-315	1	Spring	2010
		Comp. Sci.			1	Fall	2009
	Wu	Finance	90000	FIN-201	1	Spring	2010
15151	Mozart	Music	40000	MU-199	1	Spring	2010
22222	Einstein	Physics	95000	PHY-101	1	Fall	2009
32343	El Said	History	60000	HIS-351	1	Spring	2010
45565	Katz	Comp. Sci.	75000	CS-101	1	Spring	2010
45565	Katz	Comp. Sci.	75000	CS-319	1	Spring	2010
76766	Crick	Biology	72000	BIO-101	1	Summer	2009
76766	Crick	· · ·	72000	BIO-301	1	Summer	2010



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    - rename / string operations / ordering / predicates
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# The Rename Operation (1/2)

- The SQL allows renaming relations and attributes using the **as** clause: old-name **as** new-name
- E.g.,
  - select ID, name, salary/12 as monthly\_salary from instructor



# The Rename Operation (2/2)

Find the names of all instructors who have a higher salary than some instructor in 'Comp. Sci'.

#### instructor

ID	name	dept_name	salary
10101	Srinivasan	Comp. Sci.	65000
12121	Wu	Finance	90000
15151	Mozart	Music	40000
22222	Einstein	Physics	95000
32343	El Said	History	60000
22454	0.11	DI 1	07000



# The Rename Operation (2/2)

Find the names of all instructors who have a higher salary than some instructor in 'Comp. Sci'.

### instructor

ID	name	dept_name	salary
10101	Srinivasan	Comp. Sci.	65000
12121	Wu	Finance	90000
15151	Mozart	Music	40000
22222	Einstein	Physics	95000
32343	El Said	History	60000
00454		DI 1	07000

select distinct T. name from instructor as T, instructor as S where T.salary > S.salary and S.dept\_name = 'Comp. Sci.'

Keyword **as** is optional and may be omitted instructor **as** T = instructor T



# **String Operations**

- SQL includes a string-matching operator for comparisons on character strings. The operator "like" uses patterns that are described using two special characters:
  - percent (%). The % character matches any substring.
  - underscore (\_). The \_ character matches any character.
- Find the names of all instructors whose name includes the substring "dar".

select name from instructor where name like '%dar%'



# **Ordering the Display of Tuples**

List in alphabetic order the names of all instructors select distinct name from instructor order by name

We may specify desc for descending order or asc for ascending order, for each attribute; ascending order is the default.

• Example: order by name desc

- Can sort on multiple attributes
  - Example: order by dept\_name, name



# **Where Clause Predicates**

- SQL includes a **between** comparison operator
- Example: Find the names of all instructors with salary between \$90,000 and \$100,000 (that is,  $\ge$  \$90,000 and  $\le$  \$100,000)
  - select name
     from instructor
     where salary between 90000 and 100000
- Tuple comparison
  - select name, course\_id from instructor, teaches where (instructor.ID, dept\_name) = (teaches.ID, 'Biology');



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Find courses that ran in Fall 2009 or in Spring 2010

Find courses that ran in Fall 2009 and in Spring 2010

Find courses that ran in Fall 2009 but not in Spring 2010



Find courses that ran in Fall 2009 or in Spring 2010

(select course\_id from section where sem = 'Fall' and year = 2009)
union
(select course\_id from section where sem = 'Spring' and year = 2010)

Find courses that ran in Fall 2009 and in Spring 2010

Find courses that ran in Fall 2009 but not in Spring 2010



Find courses that ran in Fall 2009 or in Spring 2010

(select course\_id from section where sem = 'Fall' and year = 2009)
union
(select course\_id from section where sem = 'Spring' and year = 2010)

Find courses that ran in Fall 2009 and in Spring 2010

(select course\_id from section where sem = 'Fall' and year = 2009)
intersect
(select course\_id from section where sem = 'Spring' and year = 2010)

Find courses that ran in Fall 2009 but not in Spring 2010



Find courses that ran in Fall 2009 or in Spring 2010

(select course\_id from section where sem = 'Fall' and year = 2009)
union
(select course\_id from section where sem = 'Spring' and year = 2010)

Find courses that ran in Fall 2009 and in Spring 2010

(select course\_id from section where sem = 'Fall' and year = 2009)
intersect
(select course\_id from section where sem = 'Spring' and year = 2010)

Find courses that ran in Fall 2009 but not in Spring 2010

(select course\_id from section where sem = 'Fall' and year = 2009)
except
(select course\_id from section where sem = 'Spring' and year = 2010)



### Set operations union, intersect, and except

- Each of the above operations automatically *eliminates duplicates*
- To retain all duplicates use the corresponding multiset versions union all, intersect all and except all.



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# **Null Values**

- It is possible for tuples to have a null value, denoted by *null*, for some of their attributes
- *null* signifies an **unknown value** or that **a value does not exist**.
- The result of any arithmetic expression involving null is null
  - Example: 5 + *null* returns null
- The predicate is null can be used to check for null values.
  - Example: Find all instructors whose salary is null.

select name from instructor where salary is null


## **Null Values and Three Valued Logic**

Any comparison with *null* returns *unknown*Example: 5 < null or null <> null or null = null
Three-valued logic using the truth value *unknown*:
OR: (*unknown* or *true*) = (*unknown* or *false*) = (*unknown* or *unknown*) = (*unknown* or *unknown*) = (*false* and *unknown*) = (*false* and *unknown*) = (*unknown* and *unknown*) = (*unknown*) = (*unknown* 

NOT: (not unknown) =



### **Null Values and Three Valued Logic**

Any comparison with *null* returns *unknown* 

- Example: 5 < null or null <> null or null = null
- Three-valued logic using the truth value unknown:
  - OR: (unknown or true) = true, (unknown or false) = unknown (unknown or unknown) = unknown
  - AND: (true and unknown) = unknown, (false and unknown) = false, (unknown and unknown) = unknown
  - NOT: (not unknown) = unknown
  - "P is unknown" evaluates to true if predicate P evaluates to unknown
- Result of **where** clause predicate is treated as *false* if it evaluates to *unknown*



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    - group by / having
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## **Aggregate Functions**

These functions operate on the multiset of values of a column of a relation, and return a value

avg: average valuemin: minimum valuemax: maximum valuesum: sum of valuescount: number of values



## **Aggregate Functions (Cont.)**

Find the average salary of instructors in the Computer Science department

select avg (salary) from instructor where dept\_name= 'Comp. Sci.';

Find the total number of instructors who teach a course in the Spring 2010 semester: assume table **teaches(ID, semester, year)** 

Find the number of tuples in the course relation

select count (\*) from course;



## **Aggregate Functions (Cont.)**

Find the average salary of instructors in the Computer Science department

select avg (salary)
 from instructor
 where dept\_name= 'Comp. Sci.';

Find the total number of instructors who teach a course in the Spring 2010 semester: assume table **teaches(ID, semester, year)** 

select count (distinct *ID*)
 from *teaches* where *semester* = 'Spring' and *year* = 2010

Find the number of tuples in the course relation

select count (\*) from course;



#### **Aggregate Functions – Group By**

Find the average salary of instructors in each department

 select dept\_name, avg (salary) from instructor group by dept\_name;

ID	name	dept_name	salary
76766	Crick	Biology	72000
45565	Katz	Comp. Sci.	75000
10101	Srinivasan	Comp. Sci.	65000
83821	Brandt	Comp. Sci.	92000
98345	Kim	Elec. Eng.	80000
12121	Wu	Finance	90000
76543	Singh	Finance	80000
32343	El Said	History	60000
58583	Califieri	History	62000
15151	Mozart	Music	40000
33456	Gold	Physics	87000
22222	Einstein	Physics	95000



## **Aggregate Functions – Group By**

Find the average salary of instructors in each department

select dept\_name, avg (salary)
 from instructor
 group by dept\_name;

ID	name	dept_name	salary
76766	Crick	Biology	72000
45565	Katz	Comp. Sci.	75000
10101	Srinivasan	Comp. Sci.	65000
83821	Brandt	Comp. Sci.	92000
98345	Kim	Elec. Eng.	80000
12121	Wu	Finance	90000
76543	Singh	Finance	80000
32343	El Said	History	60000
58583	Califieri	History	62000
15151	Mozart	Music	40000
33456	Gold	Physics	87000
22222	Einstein	Physics	95000

dept_name	avg_salary
Biology	72000
Comp. Sci.	77333
Elec. Eng.	80000
Finance	85000
History	61000
Music	40000
Physics	91000



# **Aggregation (Cont.)**

 /\* erroneous query \*/ select dept\_name, ID, avg (salary) from instructor group by dept\_name;



# **Aggregation (Cont.)**

- Attributes in **select** clause outside of aggregate functions must appear in **group by** list
  - /\* erroneous query \*/ select dept\_name, ID, avg (salary) from instructor group by dept\_name;



## **Aggregate Functions – Having Clause**

Find the names and average salaries of all departments whose average salary is greater than 42000



# **Aggregate Functions – Having Clause**

Find the names and average salaries of all departments whose average salary is greater than 42000

select dept\_name, avg (salary)
from instructor
group by dept\_name
having avg (salary) > 42000;



# **Aggregate Functions – Having Clause**

Find the names and average salaries of all departments whose average salary is greater than 42000

select dept\_name, avg (salary)
from instructor
group by dept\_name
having avg (salary) > 42000;

Note: predicates in the **having** clause are applied after the formation of groups whereas predicates in the **where** clause are applied before forming groups



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  - deletion / insertion / update
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#### **Modification of the Database – Deletion**

Delete all instructors

delete from instructor

Delete all instructors from the Finance department delete from instructor where dept\_name= 'Finance';



### **Modification of the Database – Insertion**

Add a new tuple to *course* insert into *course* 

values ('CS-437', 'Database Systems', 'Comp. Sci.', 4);

or equivalently

insert into course (course\_id, title, dept\_name, credits)
values ('CS-437', 'Database Systems', 'Comp. Sci.', 4);

Add a new tuple to *student* with *tot\_creds* set to null

insert into student
values ('3003', 'Green', 'Finance', null);

# 1

## **Modification of the Database – Updates**

Increase salaries of instructors whose salary is over \$100,000 by 3%, and all others receive a 5% raise

Write two update statements:

update instructor set salary = salary \* 1.03 where salary > 100000; update instructor set salary = salary \* 1.05 where salary <= 100000;

- The order is important
- Can be done better using the case statement (next slide)



#### **Case Statement for Conditional Updates**

Same query as before but with case statement

```
update instructor
set salary = case
when salary <= 100000 then salary * 1.05
else salary * 1.03
end
```



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  - all / some / exists / not exists / derived relations / with clause



#### **Nested Subqueries**

SQL provides a mechanism for the nesting of subqueries.

A **subquery** is a **select-from-where** expression that is nested within another query.



Find courses offered in Fall 2009 and in Spring 2010

section(course\_id, semester, year)

Find courses offered in Fall 2009 but not in Spring 2010

Thursday, February 14, 2013



Find courses offered in Fall 2009 and in Spring 2010

select distinct course\_id section(course\_id, semester, year)
from section
where semester = 'Fall' and year= 2009 and semester = 'Spring'
and year= 2010;

Find courses offered in Fall 2009 but not in Spring 2010



Find courses offered in Fall 2009 and in Spring 2010



Find courses offered in Fall 2009 but not in Spring 2010



Find courses offered in Fall 2009 and in Spring 2010



Find courses offered in Fall 2009 but not in Spring 2010



Find courses offered in Fall 2009 and in Spring 2010



Find courses offered in Fall 2009 but not in Spring 2010

select distinct course\_id
from section
where semester = 'Fall' and year= 2009 and
 course\_id not in (select course\_id
 from section
 where semester = 'Spring' and year= 2010);



## **Set Comparison**

Find names of instructors with salary greater than that of some (at least one) instructor in the Biology department.

Same query using > some clause



### **Set Comparison**

Find names of instructors with salary greater than that of some (at least one) instructor in the Biology department.

select distinct T.name
from instructor as T, instructor as S
where T.salary > S.salary and S.dept name = 'Biology';

Same query using > some clause



#### **Set Comparison**

Find names of instructors with salary greater than that of some (at least one) instructor in the Biology department.

select distinct T.name
from instructor as T, instructor as S
where T.salary > S.salary and S.dept name = 'Biology';

Same query using > some clause

select name
from instructor
where salary > some (select salary
from instructor
where dept name = 'Biology');



Find the names of all instructors whose salary is greater than the salary of all instructors in the Biology department.



Find the names of all instructors whose salary is greater than the salary of all instructors in the Biology department.

select name from instructor where salary > all (select salary from instructor where dept name = 'Biology');



## **Test for Empty Relations**

- The exists construct returns the value true if the argument subquery is nonempty.
- exists  $r \Leftrightarrow r \neq \emptyset$
- **not exists**  $r \Leftrightarrow r = \emptyset$



#### **Correlation Variables**

Yet another way of specifying the query "Find all courses taught in both the Fall 2009 semester and in the Spring 2010 semester"

select course\_id
from section as S
where semester = 'Fall' and year= 2009 and
 exists (select \*
 from section as T
 where semester = 'Spring' and year= 2010
 and S.course\_id= T.course\_id);

- Correlated subquery
- Correlation name or correlation variable



#### **Not Exists**

Find all students who have taken all courses offered in the Biology department.

- Note that  $X Y = \emptyset \iff X \subseteq Y$
- Note: Cannot write this query using = all and its variants



#### **Not Exists**

Find all students who have taken all courses offered in the Biology department.

```
Note that X - Y = \emptyset \iff X \subseteq Y
```

Note: Cannot write this query using = all and its variants



#### **Derived Relations**

SQL allows a subquery expression to be used in the **from** clause

Find the average instructors' salaries of those departments where the average salary is greater than \$42,000."

select dept\_name, avg\_salary
from (select dept\_name, avg (salary) as avg\_salary
 from instructor
 group by dept\_name)
where avg\_salary > 42000;

Note that we do not need to use the **having** clause



#### With Clause

- The with clause provides a way of defining a temporary view whose definition is available only to the query in which the with clause occurs.
- Find all departments with the maximum budget

with max\_budget (value) as
 (select max(budget)
 from department)
select budget
from department, max\_budget
where department.budget = max\_budget.value;



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- Null Values
- Aggregate Functions
  - group by / having
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  - deletion / insertion / update
- Nested Subqueries
  - all / some / exists / not exists / derived relations / with clause